## Bounding the peak sidelobe level of binary sequences of all lengths

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## Abstract

Improving upon 2010 results of Alon, Litsyn, and Shpunt, it was shown in 2014 by Schmidt that asymptotically, almost all binary sequences of length n have peak sidelobe level close to  $\sqrt{2n \log n}$ . One specific result of Alon, Litsyn, and Shpunt is that if we fix  $\varepsilon > 0$ , then almost all binary sequences of length n have peak sidelobe level at most  $\sqrt{2n(\log n - (1.5 - \varepsilon) \log \log n)}$ , in the sense that the probability of not satisfying that bound approaches 0 as n approaches infinity. In this note, we prove that for all sequence lengths n > 1, there is a binary sequence of length n with peak sidelobe level at most  $\sqrt{2n(\log n - \log \log n + 0.862)}$ .

By a **binary sequence** of **length** n, we mean an n-tuple

$$A = (a_0, a_1, \ldots, a_{n-1})$$

where each  $a_j$  is +1 or -1. For  $0 \le k \le n-1$ , we define the (acyclic or aperiodic) **autocorrelations** of A by

$$c_k = \sum_{j=0}^{n-k-1} a_j a_{j+k}.$$

Informally,  $c_k$  measures how much the sequence A resembles a version of itself that has been shifted by k positions.

We let  $\mathcal{B}_n$  denote the set of all  $2^n$  binary sequences of length n. For any  $A \in \mathcal{B}_n$ , we have  $c_0 = n$ . We refer to  $c_1, \ldots, c_{n-1}$  as the **nontrivial** autocorrelations of A. An old problem, arising in communications engineering but also of interest as a stand-alone combinatorial problem, involves trying to find binary sequences in  $\mathcal{B}_n$  whose nontrivial autocorrelations are 'close' to zero in some sense.

For any  $A \in \mathcal{B}_n$ , we define the **peak sidelobe level** (PSL) of A by

$$\mu(A) = \max_{1 \le k \le n-1} |c_k|.$$

We consider A to be a 'good' sequence if  $\mu(A)$  is small. If A is a constant sequence, then trivially  $\mu(A) = n - 1$ , but very informally speaking, if A is 'random' then  $\mu(A)$  tends to be significantly smaller than O(n). Many authors have investigated upper bounds for  $\mu(A)$ . (For an excellent survey, see [3].) We might try to find upper bounds for  $\mu(A)$  that hold for some sequences  $A \in \mathcal{B}_n$ , or that hold for almost all sequences  $A \in \mathcal{B}_n$ .

To make this more precise, we turn  $\mathcal{B}_n$  into a probability space by supposing the  $a_j$  are independent Rademacher variables (i.e., random variables each equally likely to be +1 or -1). This is equivalent to assigning equal weight to each of the  $2^n$  sequences in  $\mathcal{B}_n$ , and for any function f(n), the probability that  $\mu(A) \leq f(n)$  is equal to the proportion of sequences  $A \in \mathcal{B}_n$  that satisfy  $\mu(A) \leq f(n)$ . We say  $\mu(A) \leq f(n)$  for 'almost all' sequences  $A \in \mathcal{B}_n$  if

$$\lim_{n \to \infty} \Pr[\mu(A) \le f(n)] = 1.$$

We also define

$$\mu_{\min}(n) = \min_{A \in \mathcal{B}_n} \mu(A)$$

so then if  $\mu(A) \leq f(n)$  for a nonzero proportion of sequences  $A \in \mathcal{B}_n$ , we have  $\mu_{\min}(n) \leq f(n)$ .

In 2014, Schmidt proved [7] (improving upon previous results by Alon, Litsyn & Shpunt [1], the current author [4], and Moon & Moser [5]) that if we fix  $\varepsilon > 0$ , then the probability

$$\mathbf{Pr}\Big[(\sqrt{2}-\varepsilon)\sqrt{n\log n} \le \mu(A) \le (\sqrt{2}+\varepsilon)\sqrt{n\log n}\Big]$$
(1)

approaches 1 as n approaches infinity (informally, almost all sequences  $A \in \mathcal{B}_n$  have peak sidelobe level 'close' to  $\sqrt{2n \log n}$ ). Here and throughout this article, 'log' means natural log.

Earlier, Schmidt [6] gave an explicit construction showing that for each n > 1, there is a sequence  $A \in \mathcal{B}_n$  satisfying  $\mu(A) \leq \sqrt{2n \log(2n)}$ . He also gave numerical evidence for the conjecture that his sequences satisfy  $\mu(A) = O(\sqrt{n \log \log n})$ . As pointed out in [3], several authors have conjectured that there is an infinite family of binary sequences satisfying  $\mu(A) = O(\sqrt{n})$ , but this has not been proved. In fact, the best upper bounds that have been proved to hold *either* for a positive proportion of sequences or for almost all sequences appear to be of the form  $\mu(A) = O(\sqrt{n \log n})$ .

Because of the lower bound in (1), it is not possible to prove that almost all sequences  $A \in \mathcal{B}_n$  satisfy an upper bound of the form  $\mu(A) = o(\sqrt{n \log n})$ . However, if f(n) is a certain function of n that approaches infinity more slowly than  $\log n$ , it can be shown that almost all sequences  $A \in \mathcal{B}_n$  satisfy  $\mu(A) \leq \sqrt{2n(\log n - f(n))}$ . One such result is Corollary 4.3 in [1], which shows that if we fix  $\varepsilon > 0$ , then the proportion of sequences  $A \in \mathcal{B}_n$  satisfying

$$\mu(A) > \sqrt{2n(\log n - (1.5 - \varepsilon)\log\log n)}$$

is bounded above by a multiple of  $1/(\log n)^{\varepsilon}$ , and hence approaches 0 as n approaches infinity. That is, in an asymptotic sense, almost all binary sequences of length n satisfy

$$\mu(A) \le \sqrt{2n(\log n - (1.5 - \varepsilon)\log\log n)}.$$

In this note, we prove the following, which is not as good in an asymptotic sense, but which holds for all lengths n > 1.

**Proposition.** For all n > 1, the proportion of sequences  $A \in \mathcal{B}_n$  satisfying

$$\mu(A) > \sqrt{2n(\log n - \log \log n + 0.862)}$$

is strictly less than 1. It follows that for all n > 1, we have

$$\mu_{\min}(n) \le \sqrt{2n(\log n - \log \log n + 0.862)}.$$

In the proof of the proposition, we will need the following elementary lemma.

**Lemma.** If n > 1 and K is a constant, then

$$\frac{K - \log \log n}{\log n} \ge \frac{-1}{e^{K+1}}.$$

**Proof.** Consider the function

$$f(x) = \frac{K - \log x}{x}$$

for x > 0. Using elementary calculus, we find

$$f'(x) = \frac{\log x - (K+1)}{x^2}$$

which is negative when  $0 < x < e^{K+1}$  and positive when  $x > e^{K+1}$ . It follows that for all x > 0, we have

$$f(x) \ge f(e^{K+1}) = \frac{-1}{e^{K+1}}$$

and therefore for all n > 1, we have

$$\frac{K - \log \log n}{\log n} = f(\log n) \ge \frac{-1}{e^{K+1}}.$$

## **Proof of Proposition:**

As mentioned before, we turn  $\mathcal{B}_n$  into a probability space by supposing the  $a_j$  to be independent Rademacher variables, which is equivalent to assigning equal weight to all  $2^n$  sequences in  $\mathcal{B}_n$ .

Note that the autocorrelation

$$c_k = a_0 a_k + a_1 a_{k+1} + \dots + a_{n-k-1} a_{n-1}$$

is a sum of n - k terms, each of which is  $\pm 1$ . In fact, those n - k terms are independent. (This is straightforward but not quite trivial; for a proof, see [4].) If  $1 \le k \le n - 1$ , then  $c_{n-k}$  is a sum of k independent Rademacher variables, so we can use Chernoff-type bounds (see, e.g., Corollary A.1.2 in Appendix A of [2]) to conclude that if  $\lambda > 0$ , then

$$\mathbf{Pr}\big[|c_{n-k}| > \lambda\big] < 2\exp(-\lambda^2/2k).$$

Let  $\lambda = \sqrt{2n\psi(n)}$ , where we define

$$\psi(n) = \log n - \log \log n + 0.862.$$

We then have

$$\mathbf{Pr}\Big[|c_{n-k}| > \sqrt{2n\psi(n)}\Big] < 2\exp(-n\psi(n)/k).$$

We call a sequence  $A \in \mathcal{B}_n$  'good' if  $\mu(A) \leq \sqrt{2n\psi(n)}$ , and 'bad' otherwise. Then A is bad if and only if  $|c_{n-k}| > \sqrt{2n\psi(n)}$  for some  $k = 1, \ldots, n-1$ . An overestimate for  $\mathbf{Pr}[A \text{ is bad}]$  is

$$\sum_{k=1}^{n-1} \Pr\left[ |c_{n-k}| > \sqrt{2n\psi(n)} \right] < \sum_{k=1}^{n-1} 2\exp(-n\psi(n)/k).$$

Now, consider the function

$$g(x) = 2\exp(-\psi(n)/x)$$

on the interval  $x \in [\frac{1}{n}, 1]$ . The function g(x) is an increasing function of x on that interval, so a left-endpoint Riemann sum will be an underestimate for an integral:

$$\sum_{k=1}^{n-1} g\left(\frac{k}{n}\right) \frac{1}{n} < \int_{1/n}^{1} g(x) dx$$
$$\implies \sum_{k=1}^{n-1} g\left(\frac{k}{n}\right) < n \int_{1/n}^{1} g(x) dx$$
$$\implies \sum_{k=1}^{n-1} 2 \exp(-n\psi(n)/k) < 2n \int_{1/n}^{1} \exp(-\psi(n)/x) dx$$
$$\implies \mathbf{Pr}[A \text{ is bad}] < 2n \int_{1/n}^{1} \exp(-\psi(n)/x) dx.$$

We will now perform the substitution  $u = \psi(n)/x$  on this integral. We have

$$u = \psi(n)x^{-1}$$
$$du = -\psi(n)x^{-2}dx$$
$$-(x^2/\psi(n))du = dx$$
$$x = 1/n \Rightarrow u = n\psi(n)$$
$$x = 1 \Rightarrow u = \psi(n)$$
$$x = \psi(n)/u$$
$$x^2 = (\psi(n))^2/u^2$$
$$x^2/\psi(n) = \psi(n)/u^2$$
$$dx = -(x^2/\psi(n))du = -(\psi(n)/u^2)du$$

and so the above integral becomes

$$2n \int_{1/n}^{1} \exp(-\psi(n)/x) dx = 2n \int_{n\psi(n)}^{\psi(n)} \exp(-u) \left(-\frac{\psi(n)}{u^2}\right) du$$
$$= 2n\psi(n) \int_{\psi(n)}^{n\psi(n)} \frac{1}{u^2 e^u} du.$$

That is, we have

$$\mathbf{Pr}[A \text{ is bad}] < 2n\psi(n) \int_{\psi(n)}^{n\psi(n)} \frac{1}{u^2 e^u} du.$$

Now since the function  $h(u) = 1/u^2 e^u$  decreases very rapidly, a rather crude upper bound will suffice. We have

$$\int_{\psi(n)}^{n\psi(n)} \frac{1}{u^2 e^u} du < \int_{\psi(n)}^{\infty} \frac{1}{u^2 e^u} du.$$

On the interval  $u \in (\psi(n), \infty)$ , we have  $u^2 > (\psi(n))^2$ , so we have

$$\int_{\psi(n)}^{\infty} \frac{1}{u^2 e^u} du < \frac{1}{(\psi(n))^2} \int_{\psi(n)}^{\infty} e^{-u} du = \frac{1}{(\psi(n))^2} e^{-\psi(n)}.$$

This implies that we have

$$\mathbf{Pr}[A \text{ is bad}] < 2n\psi(n) \cdot \frac{1}{(\psi(n))^2} e^{-\psi(n)} = \frac{2n}{\psi(n)e^{\psi(n)}}.$$

Now since  $\psi(n) = \log n - \log \log n + 0.862$ , we have

$$\exp(\psi(n)) = \exp(\log n) \exp(-\log \log n) \exp(0.862)$$
$$= n(\log n)^{-1} e^{K}$$

where for brevity, we write K = 0.862. We then have

$$\psi(n)e^{\psi(n)} = \left(\log n - \log\log n + K\right) \cdot n(\log n)^{-1}e^{K}$$
$$= e^{K} \left(1 + \frac{K - \log\log n}{\log n}\right)n$$

and then our lemma implies

$$\psi(n)e^{\psi(n)} \ge e^{K} \left(1 + \frac{-1}{e^{K+1}}\right)n = \left(e^{K} - \frac{1}{e}\right)n.$$

Now note that

$$e^{K} - \frac{1}{e} = e^{0.862} - \frac{1}{e} > 2.00001$$

so we have

$$\frac{2n}{\psi(n)e^{\psi(n)}} < \frac{2n}{2.00001n} = \frac{2}{2.00001} < 1$$

which completes the proof of the proposition.

## References

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